# A Database of Invariant Rings 

Gregor Kemper, Elmar Körding, Gunter Malle, B. Heinrich Matzat, Denis Vogel, and Gabor Wiese

## CONTENTS

## 1. Introduction

2. Concepts of the Database
3. An Example Session
4. Sources of Examples and Attributes Stored
5. Some Conjectures

Electronic Availability

## References


#### Abstract

We announce the creation of a database of invariant rings. This database contains a large number of invariant rings of finite groups, mostly in the modular case. It gives information on generators and structural properties of the invariant rings. The main purpose is to provide a tool for researchers in invariant theory.


## 1. INTRODUCTION

Invariant theory of finite groups is a subject which has a large variety of applications, but also displays many open questions. This applies in particular to the modular case, where the characteristic of the ground field divides the group order. Consequently, much of the recent research activity went into this area; see [Benson 1993; Smith 1997] and the references there. For a general introduction into the invariant theory of finite groups we refer the reader to the survey [Stanley 1979], or the book [Smith 1995], which gives a problem-oriented presentation.

Research in invariant theory (and, in fact, many other areas of mathematics as well) greatly benefits from the availability of examples. Examples provide a means to gain experience and understanding, to find or test conjectures, search for interesting (counter-)examples, and sometimes to prove results. In invariant theory, new algorithms and the emergence of faster computers have made it possible to study problems in a way that would be impossible by hand calculations and ad hoc methods. In fact, the computational aspects of invariant theory have recently enjoyed considerable interest in their own right (as is documented in [Sturmfels 1993] and many more recent papers such as [Derksen and Kraft 1997; Kemper 1998]). With this in mind, we have decided to assemble a collection of examples, in the form of a database, and to provide it to the public as a research tool. All computations were done in
the computer algebra system Magma [Bosma et al. 1997], which has an efficient package for invariant theory [Kemper and Steel 1999]. We used the Sun computers at the IWR in Heidelberg. Currently the database contains 5922 examples, almost all modular, and takes about 100 Mbytes of storage space. The database runs with Unix operating systems. More specifically, we have tested the database with Linux and Solaris operating systems.

The software is freely available; see Electronic Availability at the end of this article.

We ask users to quote this paper when they write articles on research which involved the database.

## 2. CONCEPTS OF THE DATABASE

## Retrieval Functions

To use the database, one cannot look at all the several thousand examples with the "naked idea". Instead, significant examples must be retrieved by systematic searches. For example, a user might want to

- see examples where Noether's degree bound [1916] is violated (i.e., the maximal degree of a generating invariant exceeds the group order $|G|$ ),
- know whether in all examples the Hilbert ideal (i.e., the ideal in the polynomial ring $K[V]$ generated by all invariants of positive degree) is generated by homogeneous elements of degree at most $|G|$,
- find the invariant ring of some particular group, or of a group which is conjugate to it.

It should be clear from these examples that there is no way to define a fixed catalogue of criteria for which users can search the database. Therefore it seemed impossible to us to implement our retrieval functions within some standard database program. In fact, the only practical way how such criteria can be formulated in a language understandable to a computer is within some computer algebra system. Moreover, a user should be able to manipulate the data retrieved from the database and not just look at it. Therefore we have decided to base our retrieval functions on the computer algebra systems Magma and Maple. There is the choice to use either one of these systems (which of course must be available). We provide access functions that take a
boolean-valued function in Magma or Maple as an argument. Users can define search criteria with such functions. After a search has been done, the examples which meet the search criterion can be loaded into Magma or Maple, respectively, for closer examination. In the following section we present an example session which shows how this works. What made it easier for us to abandon the idea of using standard database software is the fact that we are dealing with a relatively small number of items, but the data stored for each item is quite large.

## Incomplete Data

A further problem that we had to find a way to handle is the inherent difficulty of computations in invariant theory. The algorithms require the computation of Gröbner bases and the solution of large systems of linear equations [Kemper 1996; Kemper and Steel 1999]. Therefore there are examples in the database where not all information could be computed. For example, it may happen that for some invariant ring the primary invariants could be computed, but the secondary invariants were found to be out of reach. We also used an algorithm from [Hughes and Kemper 2001], which for groups of order divisible by $p:=$ char $K$ but not by $p^{2}$ calculates the Hilbert series and the depth of the invariant ring with a computational cost that is similar to the evaluation of Molien's formula. Thus for (almost) all groups in the database of order not divisible by $p^{2}$ we have the Hilbert series, depth, CohenMacaulay property, and the Gorenstein property of the invariant ring, although in many cases not even a set of primary invariants is known. We did not want to exclude such examples from the database. As a consequence, the retrieval functions have to be able to deal with incomplete information. For example, a search function supplied by a user might ask something about secondary invariants. Such a search function, when applied to a ring where the secondary invariants are not known, should not return "true" or "false", but "unknown". This feature was especially hard to implement in Magma, where there is no traperror mechanism.

## Computational Difficulty

The computational difficulty also led to some problems in the creation of the database. Usually when
one performs difficult computations on a computer, one has the computer run for a while and at some point when patience runs out, one chooses to interrupt the computation and tries a different method. Obviously this approach is not feasible for computing several thousand of examples. Instead, we implemented a scheme where different steps (or groups of steps) in the computation of each invariant ring are performed by different Magma processes which are run with a time and memory limit. If such a process terminates within the limit, it stores its results to a file for subsequent use in later steps. Otherwise, the invariant ring is transferred to a "problem queue", where it can then be worked on by ad hoc or semi-automatic methods.

## Efficient Information Transfer

When running the retrieval functions, information from the database is automatically read into Magma or Maple in order to apply the search function to the invariant rings. For reasons of efficiency it is important to transfer only that part of the information about each invariant ring into Magma or Maple which is actually needed for the evaluation of the search function. To decide what the relevant data is, one might subject the search function to a syntax analysis. Since this seemed impractical to us, we chose to implement a technique for dynamically determining the required information. More precisely, information that is found to be missing for the evaluation of the search function on some ring is reloaded for this ring, and then included into the list of necessary information for subsequent evaluations of the search function.

We believe that the specific difficulties we encountered in this project generalize to many other mathematical databases, and we hope that the concepts we developed will also be applicable in other contexts as well.

## 3. AN EXAMPLE SESSION

After the database has been downloaded, it requires a minimal amount of installation. For details see the documentation supplied with the database. Then the retrieval functions in Magma or Maple can be used. We present an example session in Magma, and remark that the usage in Maple is for the most
part analogous. We start by calling the executable InvSearch. This starts up Magma, reads in the retrieval functions, and sets up the communication with the database. In the sequel we assume some basic familiarity with Magma.
(a) As a first example, suppose we are interested in the invariants of the group $G=\mathrm{SO}_{3}\left(\mathbb{F}_{5}\right)$ in the natural representation. The chances of finding the invariants of $G$ in the database are much higher if we search for groups which are conjugate to $G$ in $\mathrm{GL}_{3}\left(\mathbb{F}_{5}\right)$, rather than only for $G$ itself. A test for this is provided by the function IsGroupConjugateTo, which is part of the retrieval functions. So we type:

```
> G := SO(3,5);
> T,F,U := SearchInvariants(func<R |
    IsGroupConjugateTo(R,G)>);
```

> T ;
[ 10077 ]

The search through the database took 75 seconds. The function SearchInvariants is called with a boolean-valued function (the "search function") as argument. This function has an invariant ring $R$ as input and returns true if $R$ is the invariant ring of a group conjugate to $G$. SearchInvariants returns three lists, $T, F$, and $U$, which stand for the invariant rings for which the search function yielded true, false, or could not be evaluated, respectively. Thus we have found exactly one invariant ring of a group conjugate to $G$. Every invariant ring is identified by a unique integer, its ExampleID. These ExampleID's are listed in T,F, and U. So far, no invariant ring has been loaded into Magma. We load the one we are interested in now, and look at some of its properties.

```
> R := RequestInvariants(T[1]);
> DegreePrimaries(R);
[ 2, 6, 20 ]
> DegreeSecondaries(R);
[ 0, 25 ]
> Hypersurface(R);
true
```

(b) Next we want to test the conjecture (Conjecture 1 below) that if $K[V]^{G}$ is Cohen-Macaulay, then Noether's degree bound holds.

```
> CM,nCM,U := SearchInvariants(func<R |
    CohenMacaulay(R)>);
```

```
> #CM,#nCM,#U;
```

333011161476

This search took 19 seconds. So we have 3330 examples of Cohen-Macaulay invariant rings, 1116 examples of non-Cohen-Macaulay rings, and 1476 examples where the Cohen-Macaulay property could not be evaluated. Now we wish to single out those examples which satisfy Noether's bound from the Cohen-Macaulay invariant rings. This can be done by giving a search range as a second argument to SearchInvariants. The minimal number $k$ such that an invariant ring $R$ can be generated by invariants of degree at most $k$ is given by the function $\operatorname{Beta}(\mathrm{R})$.

```
> NB,nNB,U := SearchInvariants(func<R |
    Beta(R) le GroupOrder(R)>,CM);
> #NB,#nNB,#U;
3105 O 225
```

Thus the conjecture could be verified in 3105 cases, and there is no counterexample.

## 4. SOURCES OF EXAMPLES AND ATTRIBUTES STORED

All finite groups with noncyclic Sylow $p$-subgroup ( $p=$ char $K$ ) have an infinite number of nonisomorphic indecomposable representations over $K$. Thus there is no way in which the representations covered in our database can reach any level of comprehensiveness, and some degree of arbitrariness is therefore unavoidable in the choice of what linear groups we included in the database. This also means that for a user it will be a matter of luck if an invariant ring he or she is interested in will be contained in the database. In order to obtain a selection of examples which is not too biased in one direction or another, we decided to take our examples from the following sources:

1. all subgroups of $\mathrm{GL}_{4}\left(\mathbb{F}_{2}\right)$,
2. all 2-subgroups of $\mathrm{GL}_{5}\left(\mathbb{F}_{2}\right)$,
3. all 3-subgroups of $\mathrm{GL}_{4}\left(\mathbb{F}_{3}\right)$,
4. all subgroups of $\mathrm{GL}_{4}\left(\mathbb{F}_{3}\right)$ which can be generated by at most two elements,
5. all the exceptional irreducible complex reflection groups in characteristic 0 , according to the classification in [Shephard and Todd 1954], apart from the generating invariants for the groups number

36 and $37\left(E_{7}\right.$ and $\left.E_{8}\right)$, which are not included in the data base because of storage problems but they can be obtained from the authors upon request,
6. miscellaneous examples that seemed of special interest to us, including some small representations of quasi-simple groups,
7. an assortment of representations up to degree 7 of groups of small order.

The groups in item 7 were produced as follows. First we used the SmallGroups library in Magma to get some groups of small order. Then for each group and each prime $p$ dividing the group order, we produced many "random" representations over $\mathbb{F}_{p^{i}}(1 \leq i \leq 3)$ by forming tensor products, symmetric powers, Jacobson radicals and other standard operations of representations we already had, and then extracting indecomposable representations from these with the Meat Axe. Since decomposable representations are also of considerable interest in invariant theory, we formed direct sums of the representations obtained in this way of total degree at most 7 .

It should also be of interest what information we store for each invariant ring. The following is a partial list of attributes that we store for an invariant ring $K[V]^{G}$, wherever they could be computed.

1. The ground field $K$,
2. the dimension of $V$,
3. generators of $G$,
4. some properties of $G$, such as the group order and whether $G$ is a $p$-group ( $p=\operatorname{char} K$ ) or a solvable group,
5. some properties of the representation $V$, such as irreduciblility, or whether $G$ acts as a (pseudo-) reflection group,
6. primary invariants,
7. secondary invariants,
8. fundamental invariants, i.e., a minimal system of generators of $K[V]^{G}$,
9. syzygies, i.e., algebraic relations between the fundamental invariants,
10. "module-syzygies", i.e., linear relations between the secondary invariants over the subalgebra generated by the primary invariants,
11. the depth of $K[V]^{G}$,

## 12. the Hilbert series,

13. the Cohen-Macaulay and Gorenstein properties, and whether $K[V]^{G}$ is a complete intersection, a hypersurface, or a polynomial ring.

## 5. SOME CONJECTURES

We conclude this note by adding a few conjectures which have all been confirmed by the database. In the following, $G \leq \mathrm{GL}(V)$ is a finite linear group in dimension $n:=\operatorname{dim}(V)$.

Conjecture 1. If $K[V]^{G}$ is Cohen-Macaulay, then Noether's degree bound holds, i.e., $K[V]^{G}$ is generated by homogeneous invariants of degrees at most $|G|$.

This conjecture generalizes the fact that Noether's degree bound holds in the nonmodular case, which was recently proved in full generality [Fleischmann 2000]. We have 3330 examples of Cohen-Macaulay invariant rings in the database. Of these, 3105 are known to satisfy Noether's bound, and for the rest generating invariants are not known. On the other hand, the database contains 133 examples that violate Noether's bound. Another generalization is contained in the following conjecture.

Conjecture 2 [Derksen and Kemper 2002, 3.7.6]. Let $I \subset K[V]$ be the "Hilbert ideal", i.e., the ideal in the polynomial ring $K[V]$ generated by all homogeneous invariants of positive degree. Then I is generated (as an ideal) by homogeneous elements of degree at most $|G|$.

Clearly Conjecture 2 holds if Noether's degree bound is satisfied. But we also verified it for all 133 examples where Noether's bound fails.

Conjecture 3 (Derksen; see [Kemper 1999]). Let $f_{1}$, $\ldots, f_{n} \in K[V]^{G}$ be primary invariants of degrees $d_{1}, \ldots, d_{n}$. Then the degrees of the (corresponding) secondary invariants are bounded from above by $d_{1}+$ $\cdots+d_{n}-n$.

Conjecture 3 was proved in the Cohen-Macaulay case by [Broer 1997]. The secondary invariants are only known for 771 of the 1116 non-Cohen-Macaulay invariant rings in the database. In all 771 examples, Conjecture 3 holds.

Conjecture 4 [Kemper 1999, Conjecture 22]. The degree of the Hilbert series $H\left(K[V]^{G}, t\right)$ (as a rational function in $\mathbb{C}(t))$ is at most $-n$.

Conjecture 4 is true in the Cohen-Macaulay case, since in this case it is equivalent to Conjecture 3. We verified the conjecture for all 1116 invariant rings in the database which are not Cohen-Macaulay.

Conjecture 5. If $K[V]^{G}$ is Cohen-Macaulay and $G \leq$ $\mathrm{SL}(V)$, then $K[V]^{G}$ is Gorenstein.

Conjecture 5 is true in the nonmodular case by a result of Watanabe [1974a; 1974b]. 1916 examples in our database satisfy the hypothesis of Conjecture 5, and all are Gorenstein. On the other hand, we have 893 examples which are Cohen-Macaulay but not Gorenstein (where the groups are not contained in $\mathrm{SL}(V)$, of course).

## ELECTRONIC AVAILABILITY

The database, with documentation and software for the retrieval of data, can be obtained from ftp:// ftp.iwr.uni-heidelberg.de/pub/kemper/DataBase/.

## REFERENCES

[Benson 1993] David J. Benson, Polynomial invariants of finite groups, London Math. Soc. Lecture Note Ser. 190, Cambridge Univ. Press, Cambridge, 1993.
[Bosma et al. 1997] Wieb Bosma, John J. Cannon, and Catherine Playoust, "The Magma algebra system, I: The user language", J. Symbolic Comput. 24 (1997), 235-265.
[Broer 1997] Abraham Broer, "Remarks on invariant theory of finite groups", preprint, Université de Montréal, Montréal, 1997.
[Derksen and Kemper 2002] Harm Derksen, Gregor Kemper, Computational invariant theory, Encyclopaedia of Math. Sciences, Springer, Berlin, 2002.
[Derksen and Kraft 1997] Harm Derksen and Hanspeter Kraft, "Constructive invariant theory", pp. 221-244 in Algèbre non commutative, groupes quantiques et invariants (Reims, 1995), Sémin. Congr. 36, Soc. Math. France, Paris, 1997.
[Fleischmann 2000] Peter Fleischmann, "The Noether bound in invariant theory of finite groups", Adv. in Math. 156 (2000), 23-32.
[Hughes and Kemper 2001] Ian Hughes and Gregor Kemper, "Symmetric powers of modular representations for groups with a Sylow subgroup of prime order", $J$. of Algebra 241:2 (2001), 759-788.
[Kemper 1996] Gregor Kemper, "Calculating invariant rings of finite groups over arbitrary fields", J. Symbolic Comput. 21 (1996), 351-366.
[Kemper 1998] Gregor Kemper, "Computational invariant theory", pp. 3-26 in The curves seminar at Queen's, vol. XII, Queen's Papers in Pure and Applied Math. 114, Queen's University, Kingston, Ont., 1998.
[Kemper 1999] Gregor Kemper, "Hilbert series and degree bounds in invariant theory", pp. 249-263 in Algorithmic algebra and number theory, edited by B. Heinrich Matzat et al., Springer, Heidelberg, 1999.
[Kemper and Steel 1999] Gregor Kemper and Allan Steel, "Some algorithms in invariant theory of finite groups", in Computational methods for representations of groups and algebras (Essen, 1997), edited by P. Dräxler et al., Prog. Math. 173, Birkhäuser, Basel, 1999.
[Noether 1916] Emmy Noether, "Der Endlichkeitssatz der Invarianten endlicher Gruppen", Math. Ann. 77 (1916), 89-92.
[Shephard and Todd 1954] G. C. Shephard and J. A. Todd, "Finite unitary reflection groups", Canad. J. Math. 6 (1954), 274-304.
[Smith 1995] Larry Smith, Polynomial invariants of finite groups, A K Peters, Wellesley (MA), 1995.
[Smith 1997] Larry Smith, "Polynomial invariants of finite groups: a survey of recent developments", Bull. Amer. Math. Soc. 34 (1997), 211-250.
[Stanley 1979] Richard P. Stanley, "Invariants of finite groups and their applications to combinatorics", Bull. Amer. Math. Soc. (N.S.) 1:3 (1979), 475-511.
[Sturmfels 1993] Bernd Sturmfels, Algorithms in invariant theory, Springer, Wien, 1993.
[Watanabe 1974a] Keiichi Watanabe, "Certain invariant subrings are Gorenstein, I", Osaka J. Math. 11 (1974), 1-8.
[Watanabe 1974b] Keiichi Watanabe, "Certain invariant subrings are Gorenstein, II", Osaka J. Math. 11 (1974), 379-388.

Gregor Kemper, IWR, Universität Heidelberg, Im Neuenheimer Feld 368, 69120 Heidelberg, Germany (Gregor.Kemper@iwr.uni-heidelberg.de)

Elmar Körding, IWR, Universität Heidelberg, Im Neuenheimer Feld 368, 69120 Heidelberg, Germany (koerding@tphys.uni-heidelberg.de)

Gunter Malle, FB Mathematik/Informatik, Universität Gh Kassel, Heinrich-Plett-Str. 40, 34132 Kassel, Germany (malle@mathematik.uni-kassel.de)
B. Heinrich Matzat, IWR, Universität Heidelberg, Im Neuenheimer Feld 368, 69120 Heidelberg, Germany (matzat@iwr.uni-heidelberg.de)

Denis Vogel, IWR, Universität Heidelberg, Im Neuenheimer Feld 368, 69120 Heidelberg, Germany (vogel@mathi.uni-heidelberg.de)
Gabor Wiese, IWR, Universität Heidelberg, Im Neuenheimer Feld 368, 69120 Heidelberg, Germany (gabor@mathphys.fsk.uni-heidelberg.de)

