# Cardinality Estimates for Piecewise Congruences of Convex Polygons

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**Abstract.** Two convex polygons  $P, P' \subseteq \mathbb{R}^2$  are congruent by dissection with respect to a given group G of transformations of  $\mathbb{R}^2$  if both can be dissected into the same finite number k of polygonal pieces  $Q_1, \ldots, Q_k$ and  $Q'_1, \ldots, Q'_k$  such that corresponding pieces  $Q_i, Q'_i$  are congruent with respect to  $G, 1 \leq i \leq k$ . In that case  $\deg_G(P, P')$  denotes the smallest k with the above property.

For the group Isom<sup>+</sup> of proper Euclidean isometries we prove two general upper estimates for  $\deg_{Isom^+}(P, P')$ , the first one in terms of the numbers of vertices and the diameters of P, P', the second one depending moreover on the radii of inscribed circles. A particular result concerns regular polygons P, P'.

For the groups  $\operatorname{Sim}^+$  and  $\operatorname{Sim}$  of proper and general similarities we give upper bounds for  $\operatorname{deg}_{\operatorname{Sim}^+}(P, P')$  and  $\operatorname{deg}_{\operatorname{Sim}}(P, P')$  in terms of the numbers of vertices.

MSC 2000: 52B45 (primary), 52B05 (secondary)

Keywords: congruence by dissection, scissors congruence, piecewise congruence, equidissectable, convex polygon, isometry, similarity, translation, number of pieces

<sup>\*</sup>This research was supported by DFG grant RI 1087/3. A part of this research has been worked out during an extended stay of the author at the Institut de Mathématiques de Jussieu, Paris.

### 1. Introduction

Given a group G of affine transformations of the Euclidean plane  $\mathbb{R}^2$ , two polygons  $P, P' \subseteq \mathbb{R}^2$  are called *congruent by dissection* (or *equidissectable*) with respect to G if there exist a number  $k \in \{1, 2, ...\}$  and dissections of P into polygons  $Q_1, \ldots, Q_k$  and of P' into polygons  $Q'_1, \ldots, Q'_k$  such that, for every  $i \in \{1, \ldots, k\}$ ,  $Q_i$  and  $Q'_i$  are congruent with respect to G. Here a polygon is meant to be a union of finitely many triangles. We say that P is dissected into  $Q_1, \ldots, Q_k$  if  $P = Q_1 \cup \cdots \cup Q_k$  and the interiors of distinct pieces  $Q_i, Q_j, i \neq j$ , are disjoint. If P and P' are equidissectable, the minimal number k admitting dissections with the above property is called the *degree* of the congruence by dissection of P and P'. This optimal number of pieces is denoted by  $\deg_G(P, P')$ .

Of course, if P and P' are congruent by dissection with respect to a subgroup H of G, then P and P' are equidissectable with respect to G, too, and

$$\deg_G(P, P') \le \deg_H(P, P')$$

The present paper is mainly devoted to degree estimates for congruences by dissection of convex polygons with respect to the group Isom of Euclidean isometries and the subgroup Isom<sup>+</sup> of proper isometries.

The classical Wallace-Bolyai-Gerwien theorem says that any two polygons of the same area are equidissectable with respect to Isom (see [4] and [10, Chapter 3] for historical remarks). The group containing all translations and all central reflections is known to be the smallest subgroup of Isom satisfying the above property (see [5, 1]). However, the question for the degree of congruences by dissection is rather open. The following theorem by Hertel seems to give the first upper estimate for deg<sub>Isom</sub>(P, P') concerning general polygons of equal area.

**Theorem 1.** ([6], Satz 2) Let  $P_m$  and  $P'_n$  be an m-gon and an n-gon of the same area whose diameters are d and d', respectively. Suppose that there exist dissections of  $P_m$  into m - 2 triangles  $T_1, \ldots, T_{m-2}$  and of  $P'_n$  into n - 2 triangles  $T'_1, \ldots, T'_{n-2}$  and define

$$c = \min\{\operatorname{diam}(T_1), \ldots, \operatorname{diam}(T_{m-2}), \operatorname{diam}(T'_1), \ldots, \operatorname{diam}(T'_{n-2})\}.$$

Then

$$\deg_{\text{Isom}}(P_m, P'_n) \le 4(m-2)(n-2)\left(\frac{\max\{d, d'\}}{c} + 2\right)^2$$

Our first goal in Section 2 is an upper estimate for  $\deg_{\text{Isom}}(P_m, P'_n)$  for arbitrary convex  $P_m$  and  $P'_n$  only depending on m, n, d, and d'. We shall see that Theorem 1 gives a bound of that kind depending on m and n like a polynomial of degree 4 and of quadratic behaviour in max  $\{\frac{d}{d'}, \frac{d'}{d}\}$ . Then we establish a stronger estimate for for  $\deg_{\text{Isom}^+}(P_m, P'_n)$  quadratic in m, n and linear in max  $\{\frac{d}{d'}, \frac{d'}{d}\}$  (see Theorem 2). An important technical step to this main result concerns the piecewise congruence of triangles (see Lemma 3).

The estimate of Theorem 2 can be improved if  $P_m$  and  $P'_n$  are known to contain inscribed circles of sufficiently large radii (see Theorem 3). In particular,

 $\deg_{\text{Isom}^+}(P_m^r, P_n^r) \leq 7(m+n-1)$  for regular polygons of the same area with m and n vertices, respectively (see Theorem 4). This improves the bound

$$\deg_{\text{Isom}}(P_m^r, P_n^r) \le (2m+4)(n+1) \quad \text{for} \quad 3 \le m < n \tag{1}$$

given by Doyen and Landuyt without proof (see [2]).

Nevertheless, our estimates do not use to be sharp in particular situations. For example, isometric congruences by dissection of regular  $P_m^r$  and  $P_n^r$  with very small m, n, say  $m, n \leq 12$ , are known to require much less than 7(m + n - 1)pieces. We refer to Theobald's frequently updated web page [9].

Let us point out that the estimates for  $\deg_{\text{Isom}^+}(P_m, P'_n)$  from Theorems 2, 3, and 4 can be realized by dissections into convex pieces.

In the last section we shall summarize new degree estimates for congruences by dissection with respect to similarities. The paper is closed with a remark about the group of translations.

We use the following notations: Given three points  $x_1, x_2, x_3 \in \mathbb{R}^2$ , the symbols  $l(x_1, x_2), x_1x_2, |x_1x_2| (= d(x_1, x_2)), \angle x_1x_2x_3$ , and  $|\angle x_1x_2x_3|$  stand for the straight line passing through  $x_1, x_2$ , the line segment between  $x_1, x_2$ , the length of that segment (which is the Euclidean distance of  $x_1, x_2$ ), the angle determined by  $x_1, x_2, x_3$ , and the size of that angle, respectively. int(A), bd(A), conv(A), diam(A) = sup{ $d(x_1, x_2) : x_1, x_2 \in A$ }, and  $\lambda(A)$  denote the interior, the boundary, the convex hull, the diameter, and the area measure, respectively, of a (measurable) subset  $A \subseteq \mathbb{R}^2$ .  $\lfloor \xi \rfloor$  and  $\lceil \xi \rceil$  are the largest lower and the smallest upper integer bound of  $\xi \in \mathbb{R}$ .

### 2. Congruence by dissection with respect to isometries

# 2.1. An estimate in terms of vertex numbers and diameters based on Theorem 1

**Corollary.** Let  $P_m$  and  $P'_n$  be convex polygons of the same area whose numbers of vertices are m and n and whose diameters are d and d', respectively. Then

$$\deg_{\text{Isom}}(P_m, P'_n) \le 4(m-2)(n-2)\left(\max\{d, d'\}\max\left\{\frac{\lfloor \frac{m}{2} \rfloor}{d}, \frac{\lfloor \frac{n}{2} \rfloor}{d'}\right\} + 2\right)^2.$$

In particular

$$\deg_{\text{Isom}}(P_m, P'_n) < mn(m+n)^2 \left( \max\left\{\frac{d}{d'}, \frac{d'}{d}\right\} \right)^2.$$

For proving this corollary, we need a lower bound for the value c from Theorem 1.

**Lemma 1.** Every convex m-gon  $P_m$  admits a dissection into m-2 triangles  $T_1, \ldots, T_{m-2}$  such that

$$\min\{\operatorname{diam}(T_1),\ldots,\operatorname{diam}(T_{m-2})\} \ge \frac{1}{\lfloor \frac{m}{2} \rfloor} \operatorname{diam}(P_m).$$

One can obtain a strict inequality if  $m \neq 3$ .

*Proof.* We proceed by induction. The case m = 3 is trivial.

In the case m = 4 let  $P_4$  have the vertices  $x_1, x_2, x_3, x_4$ . Let  $d_1 = d(x_1, x_3)$ and  $d_2 = d(x_2, x_4)$  be the lengths of the diagonals, say  $d_1 \ge d_2$ . The convexity of  $P_4$  and the triangle inequality yield

$$\operatorname{diam}(P_4) = \max_{1 \le i < j \le 4} d(x_i, x_j) < d_1 + d_2.$$

Cutting  $P_4$  along  $x_1x_3$  gives a dissection into two triangles each having a diameter of at least

$$d(x_1, x_3) = d_1 \ge \frac{d_1 + d_2}{2} > \frac{\operatorname{diam}(P_4)}{2} = \frac{1}{\lfloor \frac{4}{2} \rfloor} \operatorname{diam}(P_4).$$

Now let  $m \geq 5$ . Let  $x_1, \ldots, x_m$  be the vertices of  $P_m$  in their order along the boundary of  $P_m$ .

Case 1. diam $(P_m)$  is the length of a diagonal, say of  $x_1x_k, k \in \{3, \ldots, m-1\}$ .

Then  $x_1x_k$  dissects  $P_m$  into a k-gon  $P_{k,1} = \operatorname{conv}\{x_1, \ldots, x_k\}$  and an (m - k+2)-gon  $P_{m-k+2,2} = \operatorname{conv}\{x_k, x_{k+1}, \ldots, x_m, x_1\}$ , both with diameter  $d(x_1, x_k) = \operatorname{diam}(P_m)$ . Application of the induction hypothesis to  $P_{k,1}$  and  $P_{m-k+2,2}$  yields a dissection of  $P_m$  with the required properties.

Case 2. diam $(P_m)$  is the length of an edge of  $P_m$ , say of  $x_1x_m$ . Case 2.1.  $d(x_1, x_3) > \frac{1}{\lfloor \frac{m}{m} \rfloor} \operatorname{diam}(P_m)$ .

We cut  $P_m$  along  $x_1x_3$  into the triangle  $T = \Delta x_1x_2x_3$  and the (m-1)gon  $P_{m-1} = \operatorname{conv}\{x_3, x_4, \ldots, x_m, x_1\}$  of diameter  $d(x_1, x_m) = \operatorname{diam}(P_m)$ . Then  $\operatorname{diam}(T) \ge d(x_1, x_3) > \frac{1}{\lfloor \frac{m}{2} \rfloor} \operatorname{diam}(P_m)$ . Dissection of  $P_{m-1}$  according to the induction hypothesis gives m-3 additional triangles of sufficiently large diameters. Case 2.2.  $d(x_1, x_3) \le \frac{1}{\lfloor \frac{m}{2} \rfloor} \operatorname{diam}(P_m)$ .

Then

$$d(x_3, x_m) > d(x_1, x_m) - d(x_1, x_3) \ge \left(1 - \frac{1}{\lfloor \frac{m}{2} \rfloor}\right) \operatorname{diam}(P_m) = \frac{\lfloor \frac{m-2}{2} \rfloor}{\lfloor \frac{m}{2} \rfloor} \operatorname{diam}(P_m).$$

We split  $P_m$  along  $x_3x_m$  into the quadrilateral  $P_{4,1} = \operatorname{conv}\{x_1, x_2, x_3, x_m\}$  with  $\operatorname{diam}(P_{4,1}) = d(x_1, x_m) = \operatorname{diam}(P_m)$  and into  $P_{m-2,2} = \operatorname{conv}\{x_3, \ldots, x_m\}$  with  $\operatorname{diam}(P_{m-2,2}) \geq d(x_3, x_m) > \frac{\lfloor \frac{m-2}{2} \rfloor}{\lfloor \frac{m}{2} \rfloor} \operatorname{diam}(P_m)$ . The induction hypothesis gives dissections of  $P_{4,1}$  and of  $P_{m-2,2}$  into two and into m-4 triangles, respectively. The two pieces T of  $P_{4,1}$  satisfy

$$\operatorname{diam}(T) > \frac{1}{2}\operatorname{diam}(P_{4,1}) = \frac{1}{2}\operatorname{diam}(P_m) \ge \frac{1}{\lfloor \frac{m}{2} \rfloor}\operatorname{diam}(P_m).$$

The m-2 pieces U of  $P_{m-2,2}$  have diameters

$$\operatorname{diam}(U) \ge \frac{1}{\lfloor \frac{m-2}{2} \rfloor} \operatorname{diam}(P_{m-2,2}) > \frac{1}{\lfloor \frac{m-2}{2} \rfloor} \frac{\lfloor \frac{m-2}{2} \rfloor}{\lfloor \frac{m}{2} \rfloor} \operatorname{diam}(P_m) = \frac{1}{\lfloor \frac{m}{2} \rfloor} \operatorname{diam}(P_m).$$

This completes the proof.

*Proof of the Corollary.* We apply Theorem 1 to  $P_m$  and  $P'_n$ . According to Lemma 1 there exist dissections of  $P_m$  and  $P'_n$  such that the value c from the theorem satisfies

$$c \ge \min\left\{\frac{d}{\lfloor \frac{m}{2} \rfloor}, \frac{d'}{\lfloor \frac{n}{2} \rfloor}\right\}, \quad \text{that is} \quad \frac{1}{c} \le \max\left\{\frac{\lfloor \frac{m}{2} \rfloor}{d}, \frac{\lfloor \frac{n}{2} \rfloor}{d'}\right\}.$$
 (2)

Combining this with the statement of the theorem we obtain the first estimate of the corollary. The second one is proved by

$$\begin{aligned} \deg_{\text{Isom}}(P_m, P'_n) &\leq & 4(m-2)(n-2)\big(\max\{d, d'\}\max\{\frac{m}{2}, \frac{n}{2}\} + 2\big)^2 \\ &\leq & 4(m-2)(n-2)\big(\max\{\frac{d}{d'}, \frac{d'}{d}\}\max\{\frac{m}{2}, \frac{n}{2}\} + 2\big)^2 \\ &= & (m-2)(n-2)\big(\max\{\frac{d}{d'}, \frac{d'}{d}\}\max\{m, n\} + 4\big)^2 \\ &\leq & (m-2)(n-2)\big(\max\{m, n\} + 4\big)^2\big(\max\{\frac{d}{d'}, \frac{d'}{d}\}\big)^2 \\ &\leq & (m-2)(n-2)\big((m+n-3) + 4\big)^2\big(\max\{\frac{d}{d'}, \frac{d'}{d}\}\big)^2 \\ &= & (m-2)(m+n+1)(n-2)(m+n+1)\big(\max\{\frac{d}{d'}, \frac{d'}{d}\}\big)^2 \\ &< & m(m+n)n(m+n)\big(\max\{\frac{d}{d'}, \frac{d'}{d}\}\big)^2. \quad \Box \end{aligned}$$

We close this subsection by a claim showing that Lemma 1 is sharp. In this sense the estimate (2) is best possible. This justifies the formulation of the corollary.

**Lemma 2.** For every  $m \in \{3, 4, 5, ...\}$  and every  $\varepsilon > 0$ , there exists a convex *m*-gon  $P_m$  such that every dissection of  $P_m$  into m - 2 triangles  $T_1, ..., T_{m-2}$  satisfies

$$\min\{\operatorname{diam}(T_1),\ldots,\operatorname{diam}(T_{m-2})\} < \left(\frac{1}{\lfloor \frac{m}{2} \rfloor} + \varepsilon\right)\operatorname{diam}(P_m)$$

Sketch of the proof. We assume  $m \ge 4$ , because the case m = 3 is trivial. Let  $\delta > 0$  be small. We define  $P_m = \operatorname{conv}\{x_1, \ldots, x_m\}$  where

$$x_{i} = \begin{cases} \left( \cos\left((i-1)\delta + \delta^{2}\right), \sin\left((i-1)\delta + \delta^{2}\right) \right) & \text{if } i \text{ is odd,} \\ \left( \cos(i\delta), \sin(i\delta) \right) & \text{if } i \text{ is even.} \end{cases}$$

Then diam $(P_m) = d(x_1, x_m)$  (if  $\delta$  is sufficiently small).

Suppose that  $P_m$  is dissected into m-2 triangles  $T_1, \ldots, T_{m-2}$ . It can be shown by Euler's formula that all vertices of the triangles are vertices of  $P_m$ , too. Any such triangulation contains two different triangles each of them sharing two edges with the boundary of  $P_m$ . Only one of them can contain the long edge  $x_1x_m$ . Hence at least one of them is of the form  $T = \Delta x_i x_{i+1} x_{i+2}$  with  $1 \le i \le m-2$ . Consequently, diam $(T) = d(x_i, x_{i+2}) = d(x_1, x_3) = d(x_2, x_4)$ . We obtain

$$\frac{\operatorname{diam}(T)}{\operatorname{diam}(P_m)} = \frac{d(x_1, x_3)}{d(x_1, x_m)} \quad \text{and} \quad \lim_{\delta \downarrow 0} \frac{d(x_1, x_3)}{d(x_1, x_m)} = \lim_{\delta \downarrow 0} \frac{2\delta}{\lfloor \frac{m}{2} \rfloor 2\delta} = \frac{1}{\lfloor \frac{m}{2} \rfloor}$$

Thus  $P_m$  satisfies the claim of Lemma 2 if  $\delta$  is sufficiently small.

### 2.2. Congruence by dissection of triangles

**Lemma 3.** Let T and T' be triangles of the same area having the diameters d and d', respectively. Then

$$\deg_{\mathrm{Isom}^+}(T,T') \le 4 \left\lceil \frac{1}{2} \max\left\{ \frac{d}{d'}, \frac{d'}{d} \right\} \right\rceil + 3.$$

*Proof.* We use the method of crossposing triangle strips (see [3, Chapter 12]).

Let  $T = \Delta x_0 x_1 \hat{x}_0$  with  $d = d(x_0, x_1)$ . c denotes the centre of  $x_0 \hat{x}_0$ . Let  $\tau$  be the translation mapping  $x_0$  onto  $x_1$  and let  $\sigma$  be the central reflection with respect to c. We define  $x_i = \tau^i(x_0)$ ,  $T_i = \tau^i(T)$ ,  $\hat{x}_i = \sigma(x_i)$ , and  $\hat{T}_i = \sigma(T_i)$  for  $i \in \mathbb{Z}$ . Then the triangles  $T_i, \hat{T}_i, i \in \mathbb{Z}$ , form an infinite dissection of a strip  $\Sigma$  bounded by  $l = l(x_0, x_1)$  and  $\hat{l} = l(\hat{x}_0, \hat{x}_1)$  (see Figure 1). The sizes of the inner angles



Figure 1. The triangle strip  $\Sigma$ 

of T at  $x_0, x_1, \hat{x}_0$  are denoted by  $\alpha, \beta, \gamma$ , respectively. Since  $d(x_0, x_1) = d$  is the diameter of T, the width of  $\Sigma$  is at most  $\frac{\sqrt{3}}{2}d$ . Based on T', we introduce a strip  $\Sigma'$  and respective terms  $c', \tau', \sigma', x'_i, T'_i, \hat{x}'_i, \hat{T}'_i, \ell', \hat{\ell}', \alpha', \beta', \gamma'$  analogously. We assume  $T = \Delta x_0 x_1 \hat{x}_0$  and  $T' = \Delta x'_0 x'_1 \hat{x}'_0$  to be oriented in the same way.

Without loss of generality,  $d \leq d'$ . Then  $\alpha' \leq \alpha$  or  $\beta' \leq \beta$ , because T and T' have the same area. Again without loss of generality,  $\alpha' \leq \alpha$ .

We suppose that c' = c (which can be obtained by translating T'). Finally, we assume that the intersection  $\Sigma \cap \Sigma'$  is a parallelogram P, whose vertices  $p_1, p_2, p_3, p_4$ represent the intersections  $l \cap l'$ ,  $l \cap \hat{l}'$ ,  $\hat{l} \cap \hat{l}'$ ,  $\hat{l} \cap l'$ , respectively, such that  $d(p_2, p_3) = d'$ and  $\delta = |\angle p_1 p_2 p_3| \leq \frac{\pi}{3}$  (see Figure 2). In fact, this situation can be obtained by suitably rotating  $\Sigma'$  around c, because the width of  $\Sigma$  does not exceed  $\frac{\sqrt{3}}{2}d \leq \frac{\sqrt{3}}{2}d'$ .



Figure 2. Crossposing the strips  $\Sigma$  and  $\Sigma'$ 

The midpoints  $c_1$  and  $c_3$  of  $p_1p_2$  and  $p_3p_4$  satisfy  $d(c, c_1) = d(c, c_3) = \frac{d'}{2}$  and hence agree with the centres of  $x'_1 \hat{x}'_0$  and  $x'_0 \hat{x}'_1$ , respectively. The area of P is twice that

of T' and so twice that of T, too. This shows that  $d(p_1, p_2) = d$  and the midpoints  $c_2$  and  $c_4$  of  $p_2p_3$  and  $p_1p_4$  coincide with those of  $x_1\hat{x}_0$  and  $x_0\hat{x}_1$ , respectively.

The edges of the triangles  $T_i, \hat{T}'_i, T'_i, \hat{T}'_i$  dissect P into finitely many polygons that appear in k pairs symmetric with respect to c. Each pair consists of one member contained in  $\bigcup_{i \in \mathbb{Z}} T_i$  and one member covered by  $\bigcup_{i \in \mathbb{Z}} \hat{T}_i$ . Images of the k first members under suitable integer powers of  $\tau$  form a dissection of T (see Figure 3). Similarly, we find one element in every pair of symmetric pieces of P



Figure 3. Dissections of T and T'

such that images of these k elements under suitable integer powers of  $\tau'$  constitute a dissection of T'. This shows that  $\deg_{\text{Isom}^+}(T,T') \leq k$ .

It remains to establish an upper bound for k. We prepare this by proving

$$\alpha + \beta' + \delta \le \pi. \tag{3}$$

Among all triangles  $\Delta x_0 x_1 y$  with  $y \in \hat{l}$  and of diameter  $d(x_0, x_1) = d$  the isosceles one with  $d(x_1, y) = d(x_0, x_1) = d$  and  $|\angle x_1 x_0 y| = |\angle x_0 y x_1| = \alpha_0 \ge \frac{\pi}{3}$  maximizes the size of the inner angle at  $x_0$ , in particular  $\alpha \le \alpha_0$ . Its area coincides with that of T and can be computed by  $\frac{1}{2}d^2\sin(|\angle x_0 x_1 y|) = \frac{1}{2}d^2\sin(\pi - 2\alpha_0) = \frac{1}{2}d^2\sin 2\alpha_0$ . So

 $\alpha \le \alpha_0$ , where  $\frac{\pi}{3} \le \alpha_0 < \frac{\pi}{2}$  and  $d^2 \sin 2\alpha_0 = 2\lambda(T) = \lambda(P) = dd' \sin \delta$ .

Similarly,

$$\beta' \leq \beta'_0$$
, where  $\frac{\pi}{3} \leq \beta'_0 < \frac{\pi}{2}$  and  $d'^2 \sin 2\beta'_0 = dd' \sin \delta$ .

These admit the estimate

$$dd' \sin(\pi - \delta) = \frac{1}{2} 2dd' \sin \delta$$
  

$$\leq \frac{1}{2} \left( \frac{d \cos \alpha_0}{d' \cos \beta'_0} + \frac{d' \cos \beta'_0}{d \cos \alpha_0} \right) dd' \sin \delta$$
  

$$= \frac{1}{2} \left( \frac{d \cos \alpha_0}{d' \cos \beta'_0} d'^2 \sin 2\beta'_0 + \frac{d' \cos \beta'_0}{d \cos \alpha_0} d^2 \sin 2\alpha_0 \right)$$
  

$$= dd' \left( \cos \alpha_0 \sin \beta'_0 + \cos \beta'_0 \sin \alpha_0 \right)$$
  

$$= dd' \sin(\alpha_0 + \beta'_0).$$

Hence  $\sin(\alpha_0 + \beta'_0) \ge \sin(\pi - \delta)$  and therefore  $\alpha_0 + \beta'_0 \le \pi - \delta$ , because  $\alpha_0 + \beta'_0, \pi - \delta \in \left[\frac{\pi}{2}, \pi\right]$ . This implies (3), namely  $\alpha + \beta' + \delta \le \alpha_0 + \beta'_0 + \delta \le \pi$ .

Since  $\delta < \frac{\pi}{2}$  and  $d(x_0, c_4) = \frac{d(x_0, \hat{x}_1)}{2} \le \frac{d}{2} \le \frac{d'}{2} = d(p_1, c_4)$  in the triangle  $\triangle x_0 p_1 c_4$ , we have  $p_1 \in \bigcup_{i=0}^{\infty} x_i x_{i+1}$ . Thus  $x_0 \hat{x}_0$  meets  $p_1 p_4$  as well as  $p_2 p_3$  and

splits P into two polygons symmetric with respect to c. We denote the one that contains  $p_1p_2$  by  $P_1$ . Now k is the number of pieces of P contained in  $P_1$ .

The dissection of  $P_1$  is induced by those segments  $x_i \hat{x}_{1-i}$ ,  $x_i \hat{x}_{-i}$ ,  $x'_i \hat{x}'_{1-i}$ ,  $x'_i \hat{x}'_{-i}$ that intersect  $int(P_1)$ . Among the edges of the triangles from  $\Sigma$  these are exactly the  $x_i \hat{x}_{1-i}$ ,  $x_i \hat{x}_{-i}$  satisfying  $x_i \in x_0 p_2 \setminus \{x_0, p_2\}$ , that is,  $1 \leq i \leq i_0 = \left\lceil \frac{d(x_0, p_2)}{d} \right\rceil - 1$ .

Next we shall see that  $x'_0 \hat{x}'_0$  and  $x'_1 \hat{x}'_0$  are the only edges of triangles from  $\Sigma'$  that can intersect  $int(P_1)$ . For this it suffices to show that  $\hat{x}'_0 \in p_2 p_3 \cap P_1$ , because then all  $\hat{x}'_i x'_{1-i}$ ,  $\hat{x}'_i x'_{-i}$ ,  $i \geq 1$ , are separated from  $int(P_1)$  by  $x_0 \hat{x}_0$  and all  $\hat{x}'_i x'_{-i}$ ,  $\hat{x}'_i x'_{1-i}$ ,  $i \leq -1$ , are separated from  $int(P_1)$  by  $p_1 p_2$ . We have

$$d(c, \hat{x}'_0) = \frac{d(x'_0, \hat{x}'_0)}{2} \le \frac{d'}{2} = d(c, c_1) < d(c, p_2).$$

Hence  $\hat{x}'_0$  belongs to the open half-line  $\{p_2 + \mu(p_3 - p_2) : \mu > 0\}$ . On the other hand, the slope  $\tan \alpha$  of  $x_0 \hat{x}_0$  relative to the "horizontal" straight line l is larger than the slope  $\tan(\alpha' - \delta)$  of  $x'_0 \hat{x}'_0$ , since  $\alpha' \leq \alpha$ . This yields  $\hat{x}'_0 \in p_2 p_3 \cap P_1$ . So the dissection of  $P_1$  induced by the triangles from  $\Sigma'$  is realized by the two segments  $c\hat{x}'_0, \hat{x}'_0 c_1$  and consists of the three polygons  $P_1 \cap T'$ ,  $P_1 \cap \hat{T}'_0, P_1 \cap \hat{T}'_{-1}$ .

Now the full dissection of  $P_1$  is completed by the segments  $x_i \hat{x}_{1-i}, x_i \hat{x}_{-i}$  with  $1 \leq i \leq i_0 = \left\lceil \frac{d(x_0, p_2)}{d} \right\rceil - 1$ . First let  $1 \leq i \leq i_0 - 1$ , that is,  $x_i \in x_0 p_1 \setminus \{x_0, p_1\}$ . Then each  $x_i \hat{x}_{1-i}$  or  $x_i \hat{x}_{-i}$  passes through  $\operatorname{int}(P_1 \cap T')$  and intersects at most one of  $\operatorname{int}(P_1 \cap \hat{T}'_0)$  and  $\operatorname{int}(P_1 \cap \hat{T}'_{-1})$ . Hence subdividing the previous dissection by such a segment increases the number of pieces by at most 2. After having used all these  $2(i_0 - 1)$  segments we have at most  $3 + 4(i_0 - 1) = 4i_0 - 1$  pieces.

We show now that the two remaining segments  $x_{i_0}\hat{x}_{1-i_0}$ ,  $x_{i_0}\hat{x}_{-i_0}$  together increase the total number of pieces of  $P_1$  by at most 4. We have  $x_{i_0} \in p_1 p_2 \setminus \{p_2\}$ . If  $x_{i_0} \in p_1 c_1 \setminus \{c_1\}$  we can argue as above. In the case  $x_{i_0} = c_1$  each of  $x_{i_0}\hat{x}_{1-i_0}$  and  $x_{i_0}\hat{x}_{-i_0}$  either intersects both  $\operatorname{int}(P \cap T')$  and  $\operatorname{int}(P_1 \cap \hat{T}'_0)$ , but misses  $\operatorname{int}(P_1 \cap \hat{T}'_{-1})$ , or is collinear with  $c_1\hat{x}'_0$ , or intersects only  $\operatorname{int}(P_1 \cap \hat{T}'_{-1})$ . This yields the claim. In the final case  $x_{i_0} \in c_1 p_2 \setminus \{c_1, p_2\}$  (which is displayed in Figure 2) we compute

$$|\angle p_1 c_1 \hat{x}'_0| + |\angle p_2 x_{i_0} \hat{x}_{-i_0}| = (\beta' + \delta) + \alpha \le \pi$$

by (3). Thus  $x_{i_0}\hat{x}_{-i_0}$  meets only  $P_1 \cap \hat{T}'_{-1}$ , but misses  $\operatorname{int}(P \cap T')$  and  $\operatorname{int}(P_1 \cap \hat{T}'_0)$ . The segment  $x_{i_0}\hat{x}_{1-i_0}$  may intersect all three open polygons  $\operatorname{int}(P \cap T')$ ,  $\operatorname{int}(P_1 \cap \hat{T}'_0)$ , and  $\operatorname{int}(P_1 \cap \hat{T}'_{-1})$ . However, dissecting by both  $x_{i_0}\hat{x}_{-i_0}$  and  $x_{i_0}\hat{x}_{1-i_0}$  enlarges the number of pieces of  $P_1$  by at most 4.

The last observation shows that the total number k of pieces in  $P_1$  is bounded by  $k \leq (4i_0 - 1) + 4 = 4i_0 + 3$ . Now the proof of Lemma 3 is completed by

$$\deg_{\text{Isom}^+}(T, T') \leq k \leq 4i_0 + 3 = 4\left(\left\lceil \frac{d(x_0, p_2)}{d} \right\rceil - 1\right) + 3 \leq 4\left(\left\lceil \frac{1}{d}(d(x_0, c) + d(c, c_1) + d(c_1, p_2))\right\rceil - 1\right) + 3 \leq 4\left(\left\lceil \frac{1}{d} \left(\frac{d}{2} + \frac{d'}{2} + \frac{d}{2}\right)\right\rceil - 1\right) + 3 = 4\left\lceil \frac{d'}{2d} \right\rceil + 3 = 4\left\lceil \frac{1}{2} \max\left\{\frac{d}{d'}, \frac{d'}{d}\right\}\right\rceil + 3. \square$$

### 2.3. An improved estimate in terms of vertex numbers and diameters

**Theorem 2.** Let  $P_m$  and  $P'_n$  be convex polygons of the same area whose numbers of vertices are m and n and whose diameters are d and d', respectively. Then

$$\deg_{\mathrm{Isom}^+}(P_m, P'_n) \le (m+n-5) \left( 4 \left\lceil \max\left\{ \frac{\lfloor \frac{n}{2} \rfloor d}{d'}, \frac{\lfloor \frac{m}{2} \rfloor d'}{d} \right\} \right\rceil + 3 \right).$$

In particular

$$\deg_{\mathrm{Isom}^+}(P_m, P'_n) < 2(m+n)^2 \max\left\{\frac{d}{d'}, \frac{d'}{d}\right\}.$$

The proof is prepared by two more lemmas.

**Lemma 4.** Let the area  $\lambda(T)$  of a triangle T be represented as a sum  $\lambda(T) = \lambda_1 + \cdots + \lambda_k$  of k positive real numbers. Then one can dissect T into k triangles  $T_1, \ldots, T_k$  such that

$$\lambda(T_i) = \lambda_i$$
 and  $\operatorname{diam}(T_i) > \frac{\operatorname{diam}(T)}{2}$  for  $1 \le i \le k$ .

Proof. Let  $T = \Delta x_1 x_2 x_3$  and suppose that  $d(x_1, x_2) \geq d(x_2, x_3) \geq d(x_1, x_3)$ . We fix points  $y_1, \ldots, y_{k-1}$  of the edge  $x_1 x_3$  such that  $d(x_1, y_1) = \frac{\lambda_1}{\lambda(T)} d(x_1, x_3)$ ,  $d(y_{i-1}, y_i) = \frac{\lambda_i}{\lambda(T)} d(x_1, x_3)$  for  $2 \leq i \leq k-1$ , and  $d(y_{k-1}, x_3) = \frac{\lambda_k}{\lambda(T)} d(x_1, x_3)$ . Then T splits into  $T_1 = \Delta x_1 x_2 y_1$ ,  $T_i = \Delta y_{i-1} x_2 y_i$  for  $2 \leq i \leq k-1$ , and  $T_k = \Delta y_{k-1} x_2 x_3$ . The areas of these triangles are proportional to the lengths of their edges contained in  $x_1 x_3$ . Hence  $\lambda(T_i) = \lambda_i$ ,  $1 \leq i \leq k$ . An estimate of their diameters can be obtained by the aid of the orthogonal projection  $\pi$  onto the long edge  $x_1 x_2$ , namely

$$\operatorname{diam}(T_i) \ge d(y_{i-1}, x_2) > d(\pi(y_{i-1}), x_2) > d(\pi(x_3), x_2) \ge \frac{d(x_1, x_2)}{2} = \frac{\operatorname{diam}(T)}{2}$$

for  $2 \le i \le k$ . For  $T_1$  we even have diam $(T_1) = d(x_1, x_2) = \text{diam}(T)$ .

**Lemma 5.** Let 
$$P_m$$
 and  $P'_n$  be as in Theorem 2. Then there exist dissections of  $P_m$  into  $m + n - 5$  triangles  $T_1, \ldots, T_{m+n-5}$  and of  $P'_n$  into  $m + n - 5$  triangles  $T'_1, \ldots, T'_{m+n-5}$  such that, for  $1 \le i \le m + n - 5$ ,

$$\lambda(T_i) = \lambda(T'_i), \quad \frac{d}{2\lfloor \frac{m}{2} \rfloor} < \operatorname{diam}(T_i) \le d, \quad and \quad \frac{d'}{2\lfloor \frac{n}{2} \rfloor} < \operatorname{diam}(T'_i) \le d'.$$

Proof. By Lemma 1, there exist dissections of  $P_m$  into triangles  $S_1, \ldots, S_{m-2}$  with diam $(S_i) \geq \frac{d}{\lfloor \frac{m}{2} \rfloor}$  and of  $P'_n$  into triangles  $S'_1, \ldots, S'_{n-2}$  with diam $(S'_j) \geq \frac{d'}{\lfloor \frac{n}{2} \rfloor}$ . Let  $\mu_i = \sum_{l=1}^i \lambda(S_l), \ 0 \leq i \leq m-2$ . Then the intervals  $I_i = [\mu_{i-1}, \mu_i], \ 1 \leq i \leq m-2$ , have the lengths  $\lambda(S_i)$  and together constitute a dissection of  $[0, \lambda(P_m)]$ . Similarly, let  $\nu_j = \sum_{l=1}^j \lambda(S'_l), \ 0 \leq j \leq n-2$ . The intervals  $J_j = [\nu_{j-1}, \nu_j], \ 1 \leq j \leq n-2$ , have the lengths  $\lambda(S'_j)$  and form a dissection of  $[0, \lambda(P'_n)] = [0, \lambda(P_m)]$ . The numbers  $\mu_1, \ldots, \mu_{m-3}, \nu_1, \ldots, \nu_{n-3}$  cut  $[0, \lambda(P_m)]$  into at most m+n-5 subintervals each being completely covered by some  $I_i$  and by some  $J_j$ . Hence there exists a dissection of  $[0, \lambda(P_m)]$  into closed intervals  $K_1, \ldots, K_{m+n-5}$  of positive lengths which refines both subdivisions  $\{I_1, \ldots, I_{m-2}\}$  and  $\{J_1, \ldots, J_{n-2}\}$  simultaneously.

The interval  $I_i$  of length  $\lambda(S_i)$  splits into suitable  $K_{i_1}, \ldots, K_{i_l}$  of the lengths  $\lambda_{i_1}, \ldots, \lambda_{i_l}$ . By Lemma 4,  $S_i$  can be decomposed into triangles  $T_{i_1}, \ldots, T_{i_l}$  such that

$$\lambda(T_{i_r}) = \lambda_{i_r} \quad \text{and} \quad \operatorname{diam}(T_{i_r}) > \frac{\operatorname{diam}(S_i)}{2} \ge \frac{d}{2\lfloor \frac{m}{2} \rfloor} \quad \text{for } 1 \le r \le l.$$

This gives the dissection of  $P_m$  into  $T_1, \ldots, T_{m+n-5}$ . In the same way we can dissect  $P'_n$  into triangles  $T'_1, \ldots, T'_{m+n-5}$  with

$$\lambda(T'_i) = \lambda_i$$
 and  $\operatorname{diam}(T'_i) > \frac{d'}{2\lfloor \frac{n}{2} \rfloor}$  for  $1 \le i \le m + n - 5$ 

by partitioning  $S'_1, \ldots, S'_{n-2}$ . In particular  $\lambda(T_i) = \lambda_i = \lambda(T'_i)$ . This completes the proof.

Proof of Theorem 2. We dissect  $P_m$  and  $P'_n$  by Lemma 5. Then, for every  $i \in \{1, \ldots, m+n-5\}$ , we apply Lemma 3 to  $T_i$  and  $T'_i$ . This gives dissections of  $P_m$  and  $P'_n$  proving the first estimate. The second one can be shown as follows.

$$\begin{split} \deg_{\mathrm{Isom}^{+}}(P_{m},P_{n}') &< \qquad (m+n-5)\left(4\left(\max\left\{\frac{\lfloor\frac{n}{2}\rfloor d}{d'}, \frac{\lfloor\frac{m}{2}\rfloor d'}{d}\right\}+1\right)+3\right)\right) \\ &\leq \qquad (m+n-5)\left(4\max\left\{\frac{\frac{n}{2}d}{d'}, \frac{\frac{m}{2}d'}{d}\right\}+7\right) \\ &= \qquad (m+n-5)\left(2\max\left\{\frac{nd}{d'}, \frac{md'}{d}\right\}+7\right) \\ &\leq \qquad (m+n-5)\left(2(m+n-3)\max\left\{\frac{d}{d'}, \frac{d'}{d}\right\}+7\right) \\ &\leq \qquad (m+n-5)\left(2(m+n-3)+7\right)\max\left\{\frac{d}{d'}, \frac{d'}{d}\right\} \\ &= \qquad 2(m+n-5)\left(m+n+\frac{1}{2}\right)\max\left\{\frac{d}{d'}, \frac{d'}{d}\right\} \\ &< \qquad 2(m+n)^{2}\max\left\{\frac{d}{d'}, \frac{d'}{d}\right\}. \quad \Box \end{split}$$

# 2.4. An estimate in terms of vertex numbers, diameters, and radii of inscribed circles

**Theorem 3.** Let  $P_m$  be a convex m-gon of diameter d containing a circle of radius r and let  $P'_n$  be a convex n-gon of the same area having the diameter d' and containing a circle of radius r'. Then

$$\deg_{\operatorname{Isom}^+}(P_m, P'_n) \le (m+n-1)\left(4\left\lceil \frac{1}{2}\max\left\{\frac{d}{r'}, \frac{d'}{r}\right\}\right\rceil + 3\right).$$

In particular

$$\deg_{\text{Isom}^+}(P_m, P'_n) < (m+n) \left(2 \max\left\{\frac{d}{r'}, \frac{d'}{r}\right\} + 7\right).$$

Again the problem is reduced to piecewise congruences of triangles.

**Lemma 6.** Let  $P_m$  and  $P'_n$  be as in Theorem 3. Then there exist dissections of  $P_m$  into  $k \leq m + n - 1$  triangles  $T_1, \ldots, T_k$  and of  $P'_n$  into k triangles  $T'_1, \ldots, T'_k$  such that, for  $1 \leq i \leq k$ ,

$$\lambda(T_i) = \lambda(T'_i), \quad r < \operatorname{diam}(T_i) \le d, \quad and \quad r' < \operatorname{diam}(T'_i) \le d'.$$

Proof. Let  $x_0, \ldots, x_{m-1}$  and  $x'_0, \ldots, x'_{n-1}$  be the vertices of  $P_m$  and  $P'_n$ , respectively, ordered counterclockwise along the boundaries. Let c and c' be the midpoints of the inscribed circles of the respective polygons. We define a bijection p of the half-open interval  $[0, \lambda(P_m))$  onto  $\operatorname{bd}(P_m)$  such that the counterclockwise arc from  $x_0$  to  $p(\lambda)$  along  $\operatorname{bd}(P_m)$  together with the segments  $cx_0$  and  $cp(\lambda)$  bounds a polygon of area  $\lambda$ . Similarly, we introduce  $p': [0, \lambda(P_m)) = [0, \lambda(P'_n)] \to \operatorname{bd}(P'_n)$ .

Now let  $\{0 = p^{-1}(x_0), \dots, p^{-1}(x_{m-1})\} \cup \{0 = (p')^{-1}(x'_0), \dots, (p')^{-1}(x'_{n-1})\} = \{\lambda_0, \dots, \lambda_{k-1}\}$  be ordered such that  $0 = \lambda_0 < \dots < \lambda_{k-1}$ . Of course,  $k \le m+n-1$ . We define

$$T_{i} = \begin{cases} \operatorname{conv}\{c, p(\lambda_{i-1}), p(\lambda_{i})\}, & 1 \le i \le k-1 \\ \operatorname{conv}\{c, p(\lambda_{k-1}), x_{0}\}, & i = k, \end{cases}$$

and

$$T'_{i} = \begin{cases} \operatorname{conv}\{c', p'(\lambda_{i-1}), p'(\lambda_{i})\}, & 1 \le i \le k-1, \\ \operatorname{conv}\{c', p'(\lambda_{k-1}), x'_{0}\}, & i = k, \end{cases}$$

this way obtaining dissections of  $P_m$  into  $T_1, \ldots, T_k$  and of  $P'_n$  into  $T'_1, \ldots, T'_k$ . Then  $\lambda(T_i) = \lambda_i - \lambda_{i-1} = \lambda(T'_i)$  for  $1 \le i \le k-1$  and  $\lambda(T_k) = \lambda(P_m) - \lambda_{k-1} = \lambda(T'_k)$ . Since every  $T_i$  contains the centre c as well as at least one vertex outside the inscribed circle of radius r, the lower estimate is obvious. The upper one is trivial. The triangles  $T'_i$  behave analogously.

Now Theorem 3 can be inferred from Lemma 6 as Theorem 2 has been proved by Lemma 5.

### 2.5. An estimate for regular polygons

**Theorem 4.** Let  $P_m^r$  and  $P_n^r$  be regular polygons of the same area having m and n vertices, respectively. Then

$$\deg_{\mathrm{Isom}^+}(P_m^r, P_n^r) \le 7(m+n-1).$$

Theorem 4 is an immediate consequence of the following claim and of Lemma 3.

**Lemma 7.** Let  $P_m^r$  and  $P_n^r$  be regular polygons of area 1 having m and n vertices, respectively. Then there exist dissections of  $P_m^r$  into  $k \leq m + n - 1$  triangles  $T_1, \ldots, T_k$  and of  $P_n^r$  into k triangles  $T'_1, \ldots, T'_k$  such that, for  $1 \leq i \leq k$ ,

$$\lambda(T_i) = \lambda(T'_i) \quad and \quad \tfrac{1}{2} < \tfrac{\operatorname{diam}(T_i)}{\operatorname{diam}(T'_i)} < 2.$$

*Proof.* Simple trigonometric calculations show that the radius  $r_m$  of the largest inscribed circle, the radius  $R_m$  of the smallest circumscribed circle, and the edge length  $e_m$  of  $P_m^r$  are

$$r_m = \frac{1}{\sqrt{m \tan \frac{\pi}{m}}}, \qquad R_m = \sqrt{\frac{1 + \tan^2 \frac{\pi}{m}}{m \tan \frac{\pi}{m}}}, \qquad e_m = 2\sqrt{\frac{1}{m} \tan \frac{\pi}{m}}.$$

We assume  $3 \le m < n$  without loss of generality.

Case 1. n = 4. Then m = 3. We cut  $P_3^r$  along an axis of symmetry into  $T_1, T_2$ and  $P_4^r$  along a diagonal into  $T_1', T_2'$ . Then  $\lambda(T_i) = \lambda(T_i') = \frac{1}{2}$  and  $\frac{\operatorname{diam}(T_i)}{\operatorname{diam}(T_i')} = \frac{e_3}{2R_4} = 2^{\frac{1}{2}}3^{-\frac{1}{4}} = 1.07...$  for i = 1, 2.

Case 2.  $n \ge 5$ . Now we define dissections of  $P_m^r$  into  $T_1, \ldots, T_k, k \le m+n-1$ , and of  $P_n^r$  into  $T'_1, \ldots, T'_k$  as we did in the proof of Lemma 6. We obtain in particular

$$\lambda(T_i) = \lambda(T'_i), \quad r_m < \operatorname{diam}(T_i), \quad \text{and} \quad r_n < \operatorname{diam}(T'_i)$$

for  $1 \leq i \leq k$ . For the remaining estimate of  $\frac{\operatorname{diam}(T_i)}{\operatorname{diam}(T'_i)}$  let *i* be fixed.

Case 2.1. m = 3. The triangle  $T'_i$  is of the form  $\triangle c'y'_1y'_2$ , c' being the centre of  $P^r_n$  and  $y'_1, y'_2$  lying on a common edge of  $P^r_n$ . We can estimate diam $(T'_i) = \max\{|c'y'_1|, |c'y'_2|, |y'_1y'_2|\}$  by

$$r_5 \le r_n < \operatorname{diam}(T'_i) \le \max\{R_n, R_n, e_n\} \le \max\{R_5, e_5\} = e_5.$$
(4)

Similarly, we obtain  $T_i = \triangle cy_1y_2$ , where c is the centre of  $P_3^r$  and  $y_1, y_2$  lie on a common edge of  $P_3^r$ . The respective triangle  $T'_i$ , which has the same area as  $T_i$ , is contained in one of the n pairwise congruent triangles defined as the convex hulls of c' and an edge of  $P_n^r$ . Hence  $\lambda(T_i) = \lambda(T'_i) \leq \frac{1}{n} \leq \frac{1}{5}$ . Since  $r_3$  is the height of  $T_i$  over the edge  $y_1y_2$ , we obtain  $\frac{|y_1y_2|r_3}{2} = \lambda(T_i) \leq \frac{1}{5}$  and hence  $|y_1y_2| \leq \frac{2}{5r_3}$ . Thus diam $(T_i)$  satisfies the estimate

$$r_3 < \operatorname{diam}(T_i) = \max\{|cy_1|, |cy_2|, |y_1y_2|\} \le \max\{R_3, R_3, \frac{2}{5r_3}\} = \frac{2}{5r_3}$$

Combining this with (4) we obtain the required inequalities, namely

$$0.57... = \frac{r_3}{e_5} < \frac{\operatorname{diam}(T_i)}{\operatorname{diam}(T'_i)} < \frac{\frac{2}{5r_3}}{r_5} = 1.73..$$

Case 2.2.  $m \ge 4$ . Arguments similar to those for showing (4) give

$$r_4 \le r_m < \operatorname{diam}(T_i) \le \max\{R_m, R_m, e_m\} \le \max\{R_4, e_4\} = e_4, r_4 < r_n < \operatorname{diam}(T'_i) \le \max\{R_n, R_n, e_n\} < \max\{R_4, e_4\} = e_4.$$

Consequently,  $\frac{1}{2} = \frac{r_4}{e_4} < \frac{\operatorname{diam}(T_i)}{\operatorname{diam}(T'_i)} < \frac{e_4}{r_4} = 2$ . This completes the proof.

### 3. Remarks concerning similarities and translations

## 3.1. Congruence by dissection with respect to similarities

Any two polygons P, P' are congruent by dissection with respect to the group Sim of similarities, since P and the similar image  $\sqrt{\frac{\lambda(P)}{\lambda(P')}}P'$  of P' have the same area and hence are congruent by dissection with respect to Isom by the Wallace-Bolyai-Gerwien theorem. It is shown in [7] that any convex m-gon  $P_m$  and any convex n-gon  $P'_n$  satisfy  $\deg_{\text{Sim}}(P_m, P'_n) \leq 3(\max\{m, n\} - 2)$ . This motivates the definition

$$\deg_{Sim}(m,n) = \max\{\deg_{Sim}(P_m, P'_n) : P_m \text{ a convex } m\text{-gon}, P'_n \text{ a convex } n\text{-gon}\}$$

for  $m, n \geq 3$ . We introduce  $\deg_{\text{Sim}^+}(m, n)$  analogously,  $\text{Sim}^+$  denoting the subgroup of proper similarities. Of course,

$$\deg_{\operatorname{Sim}}(m,n) = \deg_{\operatorname{Sim}}(n,m) \le \deg_{\operatorname{Sim}^+}(m,n) = \deg_{\operatorname{Sim}^+}(n,m).$$

The above mentioned estimate from [7] now reads as

$$\deg_{\text{Sim}}(m,n) \le 3(n-2) \quad \text{for all} \quad 3 \le m \le n.$$
(5)

We could show the following.

**Theorem 5.** ([8], Theorems 5 and 6) For arbitrary  $3 \le m \le n$ ,

$$\deg_{\operatorname{Sim}^{+}}(m,n) \leq \begin{cases} \left\lfloor \frac{5m-9}{2} \right\rfloor = \left\lfloor \frac{7m+8n-27}{6} \right\rfloor & \text{if } n = m, \\ \left\lfloor \frac{7m+8n-24}{6} \right\rfloor & \text{if } n > m \end{cases}$$
(6)

and

$$\deg_{\rm Sim}(m,n) \le \begin{cases} n+3 & \text{if } m = 3, \\ m+n+\lfloor \frac{m}{3} \rfloor & \text{if } 4 \le m \le 11, \\ m+n+4 & \text{if } m \ge 12. \end{cases}$$
(7)

The estimate (6) can be realized by dissections into convex pieces, whereas our construction for (7) uses up to two non-convex simple polygons in the corresponding dissections. The bound (7) is stronger than (6) if n is sufficiently large, but can be improved by (6) for certain small m, n.

The estimate (6) is better than (5) apart from the case m = n = 3, where both bounds attain the value 3. In the latter situation the estimates are sharp. This is a consequence of the following observation.

**Lemma 8.** ([8], Corollary 2) Two triangles T, T' satisfy  $\deg_{Sim}(T, T') \leq 2$  if and only if T has an angle of size  $\alpha$  and T' has an angle of size  $\alpha'$  such that  $\alpha = \alpha'$  or  $\alpha + \alpha' = \pi$ .

#### 3.2. Congruence by dissection with respect to translations

Two polygons of the same area are not necessarily equidissectable with respect to the group Trans of translations. Since directions are invariant under translations, congruence by dissection with respect to Trans depends on the directional behaviour of the corresponding polygons, too. We recall a necessary and sufficient condition from [5].

Let  $\mathbb{S}^1$  be the set of all vectors of length 1 in  $\mathbb{R}^2$ . Given a (not necessarily convex) polygon P and a direction  $x \in \mathbb{S}^1$ , l(P, x) is to denote the sum of the lengths of all edges of P having the outer normal vector x. If no such edges exist, we put l(P, x) = 0. Now let  $L_x(P) = l(P, x) - l(P, -x)$ . It is shown in [5] that two polygons P, P' are congruent by dissection with respect to Trans if and only if they have the same area and  $L_x(P) = L_x(P')$  for all  $x \in \mathbb{S}^1$ .

Suppose from now on that P, P' are arbitrary convex polygons congruent by dissection under translations. In contrast with Theorem 2, a general upper estimate for deg<sub>Trans</sub>(P, P') cannot be given in terms of the vertex numbers and the diameters of P and P'. For example, the two rectangles  $R(m) = [0, m] \times [0, 1]$  and  $R'(m) = [0, 1] \times [0, m]$  satisfy deg<sub>Trans</sub>(R(m), R'(m)) = m for every  $m \in \{1, 2, \ldots\}$ , even though their diameters agree. The last observation reveals a more general principle.

Given a bounded non-empty set  $A \subseteq \mathbb{R}^2$  and a vector  $x \in \mathbb{S}^1$ , the *width* of A in direction x is given by

$$w(A, x) = \sup\{\langle a_1 - a_2, x \rangle : a_1, a_2 \in A\},\$$

where  $\langle \cdot, \cdot \rangle$  stands for the scalar product. Of course, this functional is monotone and invariant under translations, that is,  $w(A, x) \leq w(B, x)$  if  $A \subseteq B \subseteq \mathbb{R}^2$  and w(A, x) = w(A + t, x) for all  $t \in \mathbb{R}^2$ .

Since P can be dissected into  $\deg_{\operatorname{Trans}}(P, P')$  translates  $Q_i$  of pieces of P', there is a convex set of width w(P, x) that can be covered by  $\deg_{\operatorname{Trans}}(P, P')$  sets of width  $w(Q_i, x) \leq w(P', x)$ . This yields  $\deg_{\operatorname{Trans}}(P, P') \geq \left\lceil \frac{w(P, x)}{w(P', x)} \right\rceil$  and, by symmetry,

$$\deg_{\operatorname{Trans}}(P,P') \ge \sup\left\{ \left\lceil \max\left\{\frac{w(P,x)}{w(P',x)}, \frac{w(P',x)}{w(P,x)}\right\} \right\rceil : x \in \mathbb{S}^1 \right\}$$

as a general lower estimate for  $\deg_{\text{Trans}}(P, P')$ .

We denote the right hand side of the last formula by W(P, P'). Can one find a general upper bound for deg<sub>Trans</sub>(P, P') only in terms of W(P, P') and the vertex numbers of P and P'?

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Zbl 0569.43001

Received January 24, 2008